

Double cats as pseudomonads

w app. double profunctors

Plan

1. Review 1D

- categories \approx monads
- profunctors \approx bimodules
- functors / cofunctors \approx certain special cases

2. Extend definitions to 2D

- PsDb|Cat (recover notions known in the literature)
- PsDb|Cof (introduced M. 2023)
- PsDb|Prof (WIP ju. N. Gambino)

3. Grothendieck construction

profunctors

$$A^{\text{op}} \times B \rightarrow \text{set}$$

double
profunctors

???

4. Strictification

• double
functors

(Grandis, Paré 1999)
(Campbell 2019)

• double
cofunctors

(M. 2023)

• double
profunctors

(M. 2024.)

Part 1

Review

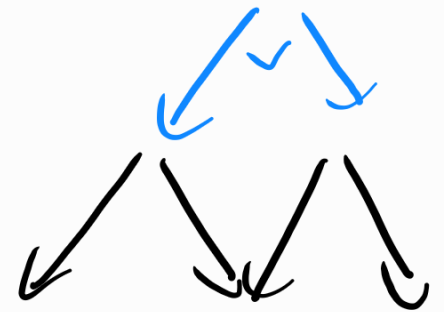
ID

case

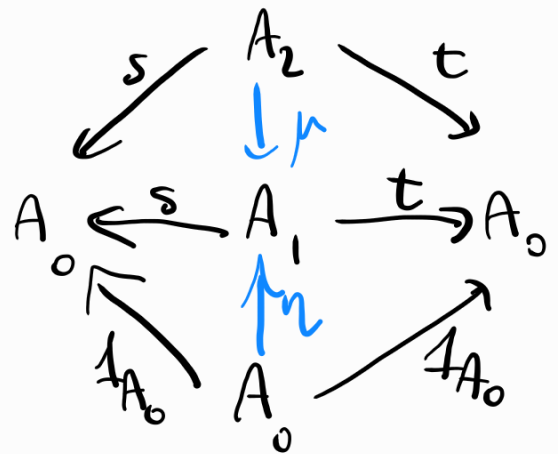
Bénabou 1967 / 1973

Span

Composition
is via pullback



monads
are categories



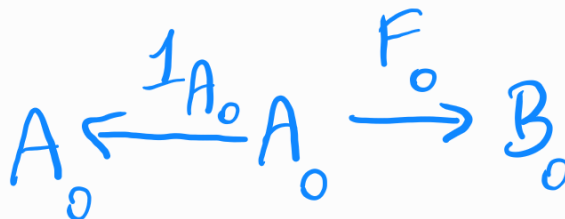
Street 1972

Formal Theory of Monads

Morphisms
of monads



cofunctors



functors

Co-morphisms
of monads



2-cells



equalities

Lack-Street 2002

KL & EM

(co)completions

More general 2-cells which correspond to natural transformations (& 2-cells bw cofunctors)

Garner Shulman 2016

Joyal Gambino 2017

Bimod (\mathcal{K})

Bimodules of monads \longleftrightarrow profunctors between cats

Morphisms of bimodules \longleftrightarrow Natural transf. between profunctors

Bimodules

$$A_0 \xleftarrow{P_1} P_1 \xrightarrow{P_0} B_0$$

P_1, α, γ
 heteromorphisms

$$\begin{array}{ccc}
 & P_1 & \\
 & \uparrow & \\
 A_0 & \xrightarrow{P_1} & B_0 \\
 \downarrow A & \nearrow P_0 & \\
 A_0 & & P_1
 \end{array}$$

$\uparrow P_1$

$$\begin{array}{ccc}
 & P_1 & \\
 & \downarrow & \\
 P_1 \times_{A_0} A_1 & \xrightarrow{P_0} & P_1
 \end{array}$$

$$\begin{array}{ccc}
 A_0 & \xrightarrow{P_1} & B_0 \\
 \searrow P_0 & & \downarrow B \\
 & & B_0 \\
 P_1 & &
 \end{array}$$

$\downarrow P_0$

$$\begin{array}{ccc}
 & P_1 & \\
 & \downarrow & \\
 B_1 \times_{B_0} P_1 & \xrightarrow{A_0} & P_1
 \end{array}$$

$$A^{op} \times B \xrightarrow{P} \text{Set}$$

algebras for the distributive law between $B^0(-)$ & $(-)^0 A_0$,
 on $\text{Span}(\text{Set}) (A_0, B_0)$

Composition of bimodules

$$A \xrightarrow{P} B \xrightarrow{Q} C$$

... is given by the **reflexive coequaliser**

$$Q \circ B \circ P \begin{array}{c} \xrightarrow{j^Q \circ 1} \\ \xrightarrow{1 \circ \lambda^P} \end{array} Q \circ P \longrightarrow Q \circ_B P$$

Cofunctors and functors

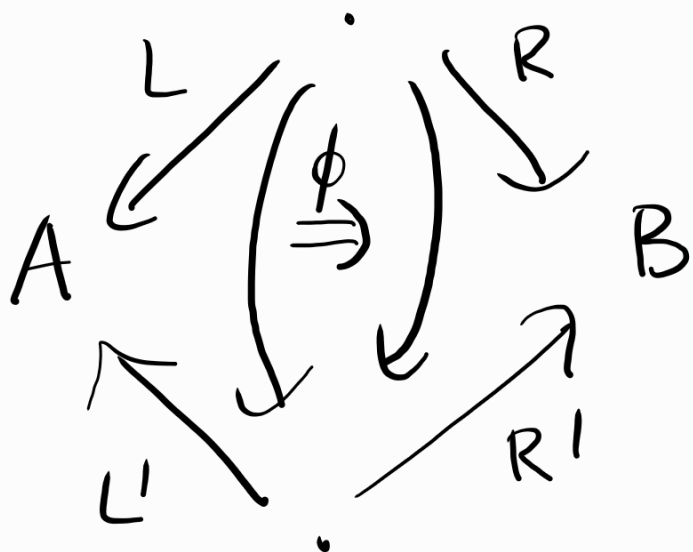
j free on $A_0 = A_0 \xrightarrow{F_0} B_0$

λ free on $A_0 = A_0 \xrightarrow{F_0} B_0$

Part 2

Tricategory

Span (Cat)

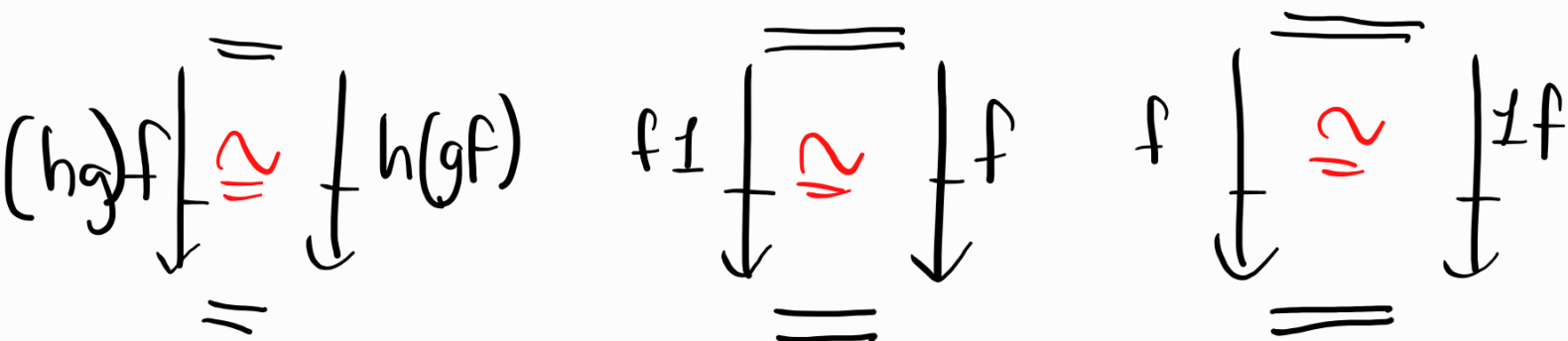


$$R'\phi = R$$
$$L'\phi = L$$

Gambino, Lobbia 2021

Formal Theory of Pseudomonads pt 1.

pseudomonads in Span (Cat) \iff (loosely pseudo) double cat's



Co-morphisms

$$A_0 \xleftarrow{1_{A_0}} A_0 \xrightarrow{F_0} B_0$$



(loosely pseudo)

double functors

2-cells



(tightly natural,
loosely pseudonatural.)

double n.t w/
identity components.

3-cells



equalities

M. 2023

Formal Theory of
Pseudomonads pt 2.

More general
2-cells



loose ps. nat
transformations

More general
3-cells



double
modifications

Morphisms of pseudomonads give a notion of (loosely pseudo) double cofunctors.

$$A_0 \xrightarrow[\text{functor}]{F_0} B_0$$

$$\begin{array}{ccc}
 X \xrightarrow{x} X' & & F_0 X \xrightarrow{F_0 x} F_0 X' \\
 \downarrow F_x f & \leftarrow & \downarrow f \\
 F_x \beta & & \Downarrow \beta \\
 Y \xrightarrow{y} Y' & & Y \xrightarrow{y} Y' \\
 \downarrow F_{x'} f' & & \downarrow f'
 \end{array}$$

respecting loose composition +

identities up to coherent inv.

double cells.

General 2-cells/3-cells give a notion of double pseudonatural transformation / double comodification

$$\begin{array}{ccc}
 \begin{array}{ccc}
 X & \xrightarrow{x} & Y \\
 \phi_x \downarrow & \Downarrow \phi_x & \downarrow \phi_y \\
 X' & \xrightarrow{x'} & Y'
 \end{array} & &
 \begin{array}{ccc}
 X & \xrightarrow{F_x f} & Y \\
 \phi_x \downarrow & \sim_{\phi_{x,f}} & \downarrow \phi_y \\
 X' & \xrightarrow{G_{x',f}} & Y'
 \end{array}
 \end{array}$$

Thm (M. 2023)

These data assemble into a tricategory PSDbICof .

N.B. New even for bicategories.

What about double profunctors?

Problem:

$$\text{Cat}/B \xrightarrow{F^*} \text{Cat}/A$$

does not preserve
coequalisers, but

Propⁿ (M. 2024)

F^* does preserve
codescent objects.

WIP j.w N. Gambino :

If \mathcal{K} is a Gray-cat
whose homs have
codescent objects separately
preserved by composition
then there is a tricategory

$\text{Bimod}(\mathcal{K})$.

objects : pseudomonads,

homs : pseudoalgebras for
pseudo-distributive law,

Composition : codescent
objects.

Part 3 Grothendieck construction.

$$A_0 \longleftarrow P_1 \longrightarrow B_0$$

(loose) heteromorphisms,

hetero-doublecells

$$\begin{array}{ccc} X & \xrightarrow{P} & Y \\ f \downarrow & \Downarrow \tilde{P} & \downarrow g \\ X' & \xrightarrow{P'} & Y' \end{array}$$

$$P_1 \longrightarrow A_0 \times B_0$$

$$A_0 \times B_0 \rightsquigarrow \text{Prof}$$

normal
lax

P_1 has actions on either side by A & B :

$$\begin{array}{ccc}
 P_1 \times_{A_0} A_1 & \xrightarrow{r} & P_1 \\
 B_1 \times_{B_0} P_1 & \xrightarrow{l} & P_1
 \end{array}
 \left. \vphantom{\begin{array}{ccc} P_1 \times_{A_0} A_1 & \xrightarrow{r} & P_1 \\ B_1 \times_{B_0} P_1 & \xrightarrow{l} & P_1 \end{array}} \right\} \text{functors}$$

$$\left(\begin{array}{c} X \\ \downarrow f \\ X' \end{array}, Y \right) \longmapsto P(f, Y) \begin{array}{c} \uparrow \\ P(X, Y) \\ \downarrow \\ P(X', Y) \end{array} \text{functor}$$

$$\left(X, \begin{array}{c} Y \\ \downarrow g \\ Y' \end{array} \right) \longmapsto P(X, g) \begin{array}{c} \downarrow \\ P(X, Y) \\ \uparrow \\ P(X, Y') \end{array} \text{functor}$$

The 2-cells of spans r & l underlie pseudoalgebras for

$$(-) \circ A, \quad B \circ (-).$$

$$\begin{array}{ccc}
 P_1 \times A_1 \times A_1 & \xrightarrow{r \times 1} & P_1 \times A_1 \\
 \downarrow 1 \times \mu & \cong \chi & \downarrow r \\
 P_1 \times A_1 & \xrightarrow{r} & P_1
 \end{array}
 \qquad
 \begin{array}{ccc}
 P_1 & \xrightarrow{1 \times \eta} & P_1 \times A_1 \\
 \downarrow 1_{P_1} & \cong \cup & \downarrow r \\
 P_1 & & P_1
 \end{array}$$

+ pseudoalgebra axioms

similar for l .

Thus we have pseudofunctors

$$A_{\text{loose}}^{\text{op}} \xrightarrow{P(-, Y)} \text{Cat}, \quad B_{\text{loose}} \xrightarrow{P(X, -)} \text{Cat}$$

The bimodularity condition is mediated by an invertible natural transformation.

$$\begin{array}{ccc}
 & B_1 \times P_1 \times A_1 & \\
 \swarrow \scriptstyle{l \times 1} & & \searrow \scriptstyle{1 \times r} \\
 P_1 \times A_1 & \cong & B_1 \times P_1 \\
 \searrow \scriptstyle{r} & & \swarrow \scriptstyle{l} \\
 & P_1 &
 \end{array}$$

Thm (M. 2024)

These data comprise a double functor

$$\left(A^{\text{op}} \times B \right)^T \rightsquigarrow \underline{\text{Prof}}$$

that is (loosely) normal lax,

and (tightly) pseudo.

N.B. This formulation is hard to guess without the pseudomanad perspective on double categories!

Composition of 1D profunctors

$$\left\{ (x \xrightarrow{p} y, y \xrightarrow{q} z) \right\}$$

$$\left\langle (p, q) \sim (p', q') \mid \begin{array}{c} y \\ \exists \downarrow \\ y' \end{array} \quad (y \xrightarrow{p}, q) = (p', q' \xrightarrow{y}) \right\rangle$$

Composition of double profunctors

$$\begin{array}{ccccc} X & \xrightarrow{y p} & Y' & \xrightarrow{q} & Z \\ \parallel & & \cong & & \parallel \\ X & \xrightarrow{p} & Y & \xrightarrow{q y} & Z \end{array}$$

Part 4: Strictification

(Grandis,
Paré
1999)

$$A \begin{cases} \cong \\ \downarrow \\ \text{st}_2(A) \end{cases}$$

weak in the
loose direction

strict in
both directions

$$\text{DblCat} \begin{array}{c} \xleftarrow{\text{st}_2(-)} \\ \perp \\ \xrightarrow{\text{Inc}} \end{array} \text{PsDblCat}$$

(Campbell 2019)

Th^M (Grandis / Paré
M., Gambino / M.)

Let

A, B be strict double
categories & $A \xrightarrow{P} B$
be a double profunctor,

1. $(A_{\text{loose}})_\perp$ free on a graph, P a double pseudofunctor

$\Rightarrow P \cong P'$ a strict double functor (Grandis' Parè)

2. $(B_{\text{loose}})_\perp$ " " "

" P a double pseudocofunctor

$\Rightarrow P \cong P'$ a strict

double cofunctor (M. 2023)

Th^m (M. 2024)

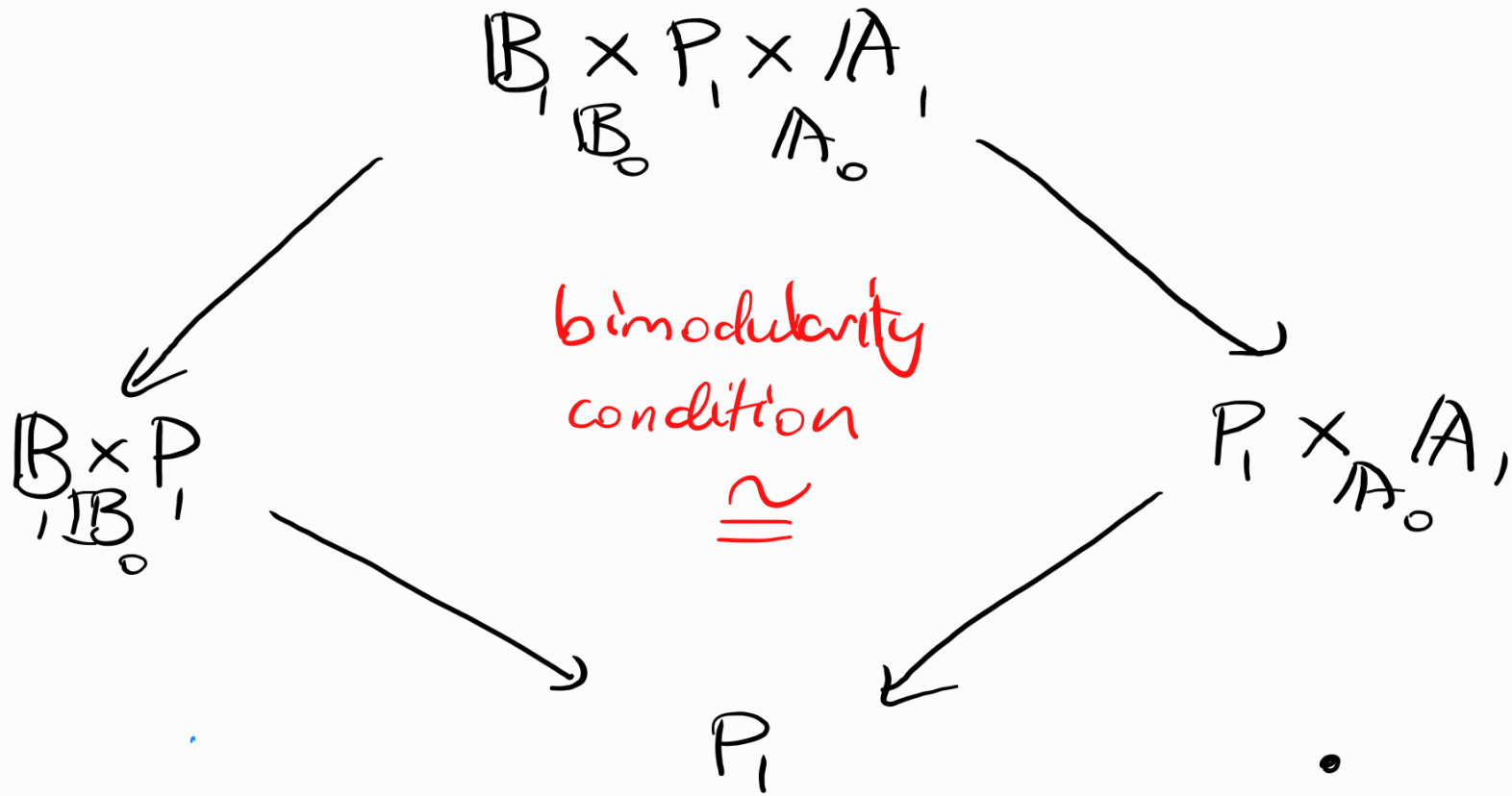
$(A_{\text{loose}})_1$, $(B_{\text{loose}})_1$ are
both free on graphs,

$A \xrightarrow{P} B$ a double
profunctor

$\Rightarrow P \cong P'$ where

$$\begin{array}{ccc} P_1 \times_{A_0} A & \xrightarrow{r} & P_1 \\ B_1 \times_{B_0} P_1 & \xrightarrow{l} & P \end{array} \left. \vphantom{\begin{array}{ccc} P_1 \times_{A_0} A & \xrightarrow{r} & P_1 \\ B_1 \times_{B_0} P_1 & \xrightarrow{l} & P \end{array}} \right\} \text{strict algebras}$$

... but still have a natural iso. mediating the usual



$$\begin{array}{c}
 (X \xrightarrow{x} X' \xrightarrow{p} Y') \xrightarrow{y} Y \\
 \parallel \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \parallel \\
 X \xrightarrow{x} (X' \xrightarrow{p} Y' \xrightarrow{y} Y)
 \end{array}$$

\cong

These semi-strict double profunctors correspond to double functors

$$\left(A \overset{\text{loose}}{\text{op}} \times B \right)^T \rightsquigarrow \text{Prof}$$

loose: normal lax

tight: separately strict,
jointly pseudo.

A strict $\Rightarrow A \xrightarrow{1_A} A$
Strict

$A \xrightarrow{P} B \xrightarrow{Q} C$

A, B, C strict

P, Q semi-strict

$\Rightarrow Q \circ P$ semi-strict.