

# COFUNCTORS, MONOIDS, AND SPLIT EPIMORPHISMS

## AUSTRALIAN CATEGORY SEMINAR

### 28 OCTOBER 2020

#### §0. MOTIVATION

A **category** is like a **monoid** with several objects.

But what is the correct generalisation of **monoid homomorphisms**?

Obvious answer: A **functor** between categories.

Thus monoids and homomorphisms form a full subcategory of  $\text{Cat}$ :

$$\boxed{\text{Mon}} \xrightarrow{\text{f.f.}} \boxed{\text{Cat}}$$

Remark: this inclusion functor has a **left adjoint**, which maps a small category to the free monoid on its set of morphisms modulo some relations on identities and composites; however no right adjoint exists as initial objects are not preserved.

Less obvious answer: A **cofunctor** between categories.

The category of small categories and cofunctors also has  $\text{Mon}$  as a full subcategory:

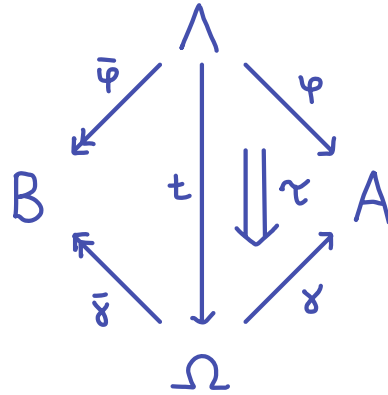
$$\boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\quad} \\ \sqcap \end{array} \boxed{\text{Cof}}$$

This inclusion functor has a **right adjoint**, which maps a small category to its **monoid of admissible sections**.

## Outline of the talk:

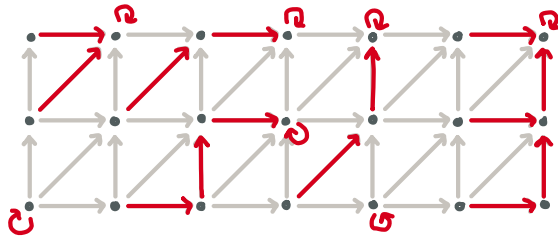
### §1. The 2-category of categories and cofunctors

- What is a cofunctor?
- Examples
- Cofunctors as spans of functors
- Morphisms of cofunctors



### §2. A right adjoint: the monoid of admissible sections

- An explicit description (vector fields on categories)



- Examples
- Characterisation as a hom category
- An internal perspective

### §3. Lenses and split epimorphisms of monoids

- Motivation: generalising semidirect products of groups
- Schreier split epimorphisms between monoids

$$K[f] \begin{array}{c} \xleftarrow{\quad q \quad} \\ \xrightarrow{\quad k \quad} \end{array} A \begin{array}{c} \xleftarrow{\quad \varphi \quad} \\ \xrightarrow{\quad f \quad} \end{array} B$$

- Lenses and split opfibrations between monoids

## §1. The 2-category of categories and cofunctors

Definition (Higgins, Mackenzie ; Aguiar): A **cofunctor**  $\varphi: B \rightarrow A$  between small categories consists of a pair of functions,

$$\begin{array}{ccc} A_0 & \xrightarrow{\varphi_0} & B_0 \\ a & \longmapsto & \varphi_0 a \end{array} \quad \text{assignment on objects}$$

$$\begin{array}{ccc} A_0 \times_{B_0} B_1 & \xrightarrow{\varphi_1} & A_1 \\ (a, \varphi_0 a \xrightarrow{u} b) & \longmapsto & \varphi_1(a, u): a \rightarrow a' \end{array} \quad \text{assignment on morphisms}$$

subject to the following axioms:

$$(1) \quad \varphi_0 \text{cod}(\varphi_1(a, u)) = \text{cod}(u)$$

$$(2) \quad \varphi_1(a, 1_{\varphi_0 a}) = 1_a$$

$$(3) \quad \varphi_1(a, v \circ u) = \varphi_1(a', v) \circ \varphi_1(a, u) \quad \text{where } a' = \text{cod } \varphi_1(a, u)$$

$$\begin{array}{ccc} A & & a \xrightarrow{\varphi_1(a, u)} a' \\ \uparrow \varphi & & \vdots \quad \quad \quad \vdots \\ B & & \varphi_0 a \xrightarrow{u} b = \varphi_0 a' \end{array}$$

"A cofunctor is a lifting of morphisms against objects, which respects identities and composition"

Given cofunctors  $\gamma: C \rightarrow B$  and  $\varphi: B \rightarrow A$ , we can take their composite in the following way:

$$\begin{array}{ccc}
 A & & \\
 \uparrow \varphi & & \\
 B & & \\
 \uparrow \gamma & & \\
 C & & 
 \end{array}
 \begin{array}{ccc}
 a & \xrightarrow{\varphi_1(a, \gamma_1(\varphi_0 a, u))} & a' \\
 \vdots & & \vdots \\
 \varphi_0 a & \xrightarrow{\gamma_1(\varphi_0 a, u)} & b \\
 \vdots & & \vdots \\
 \gamma_0 \varphi_0 a & \xrightarrow{u} & c
 \end{array}$$

This yields a category  $\mathbf{Cof}$  of categories and cofunctors.

### Examples:

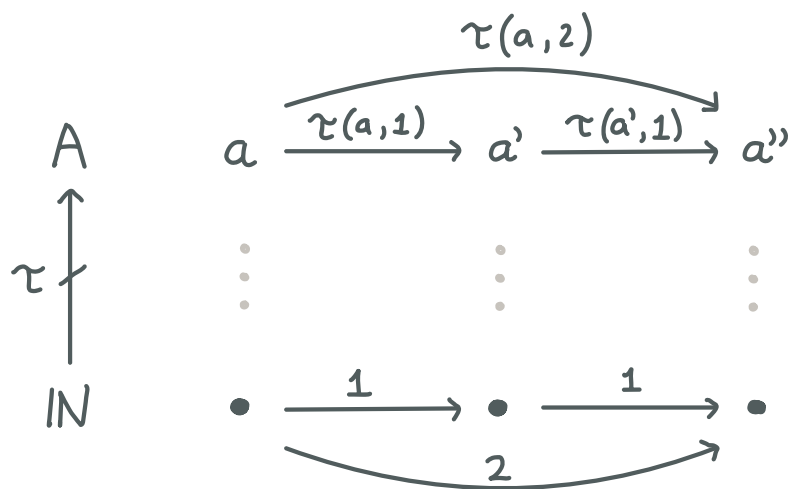
- Every **monoid homomorphism**  $B \xrightarrow{\varphi} A$  yields a cofunctor:

$$\begin{array}{ccc}
 A & & \\
 \uparrow \varphi & & \\
 B & & 
 \end{array}
 \begin{array}{ccc}
 \bullet & \xrightarrow{\varphi(m)} & \bullet \\
 \vdots & & \vdots \\
 \bullet & \xrightarrow{m} & \bullet
 \end{array}$$

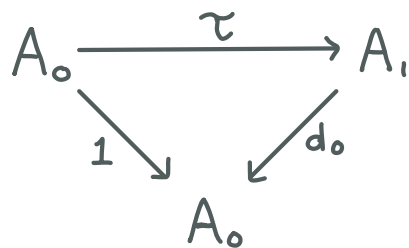
This gives a fully faithful functor  $\mathbf{Mon} \longrightarrow \mathbf{Cof}$

- More generally, every **bijective-on-objects** functor  $B \rightarrow A$  yields a cofunctor  $B \dashrightarrow A$ .
- Every **discrete opfibration**  $A \rightarrow B$  yields a cofunctor  $B \dashrightarrow A$ .

- Every **split opfibration** has an underlying cofunctor given by the splitting.
- More generally, every **delta lens** has an underlying cofunctor.
- Let  $\mathbb{N}$  denote the monoid of natural numbers under addition. A cofunctor  $\mathbb{N} \xrightarrow{\tau} A$  is the same as a choice of morphism out of every object in  $A$ :



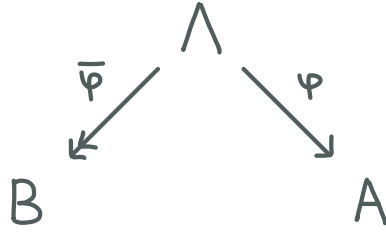
Internally this is the same as a **section** of the domain map:



- For every category  $A$ , there is a unique cofunctor  $\mathbb{1} \xrightarrow{\tau} A$  from the category with a single identity arrow. That is,  $\mathbb{1}$  is the **initial object** in  $\text{Cof}$ .

Proposition: Every cofunctor  $B \xrightarrow{\varphi} A$  may be represented as a span of functors,

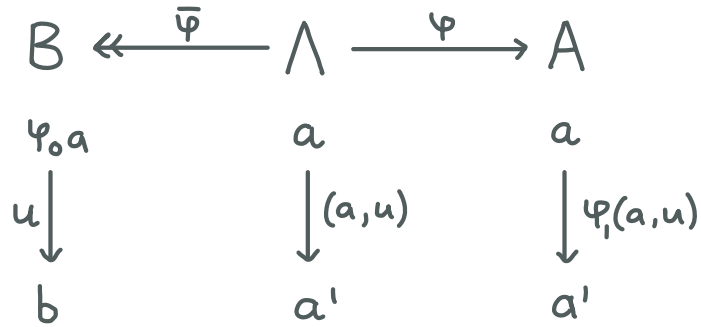
use  $\longrightarrow$  for discrete opfibrations



where  $\bar{\varphi}$  is a **discrete opfibration** and  $\varphi$  is **identity-on-objects**.

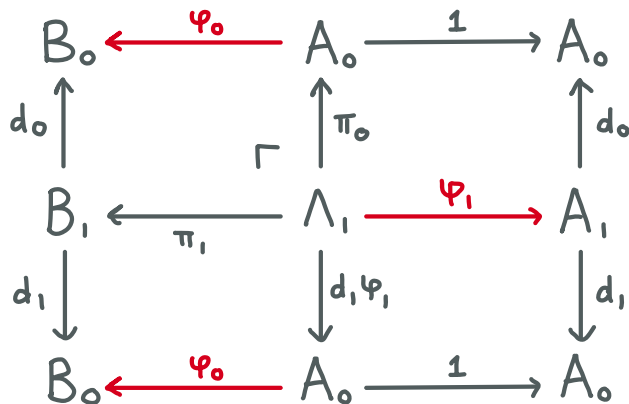
This representation extends to a functor  $\text{Cof} \longrightarrow \text{Span}_{\text{iso}}(\text{Cat})$ , and the pair  $(\text{DOpf}^{\text{op}}, \text{Bij})$  is an **orthogonal factorisation system** on  $\text{Cof}$ .

PROOF (SKETCH): Given  $\varphi: B \rightarrow A$ , let  $\Lambda$  be the category with the same objects as  $A$  and morphisms given by formal pairs  $(a, u: \varphi_0 a \rightarrow b)$ . Then we have:



where  $a' = \text{cod } \varphi_1(a, u)$ .

Internally this span of functors is given by the diagram:

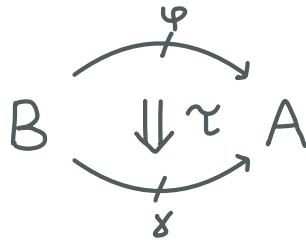


where

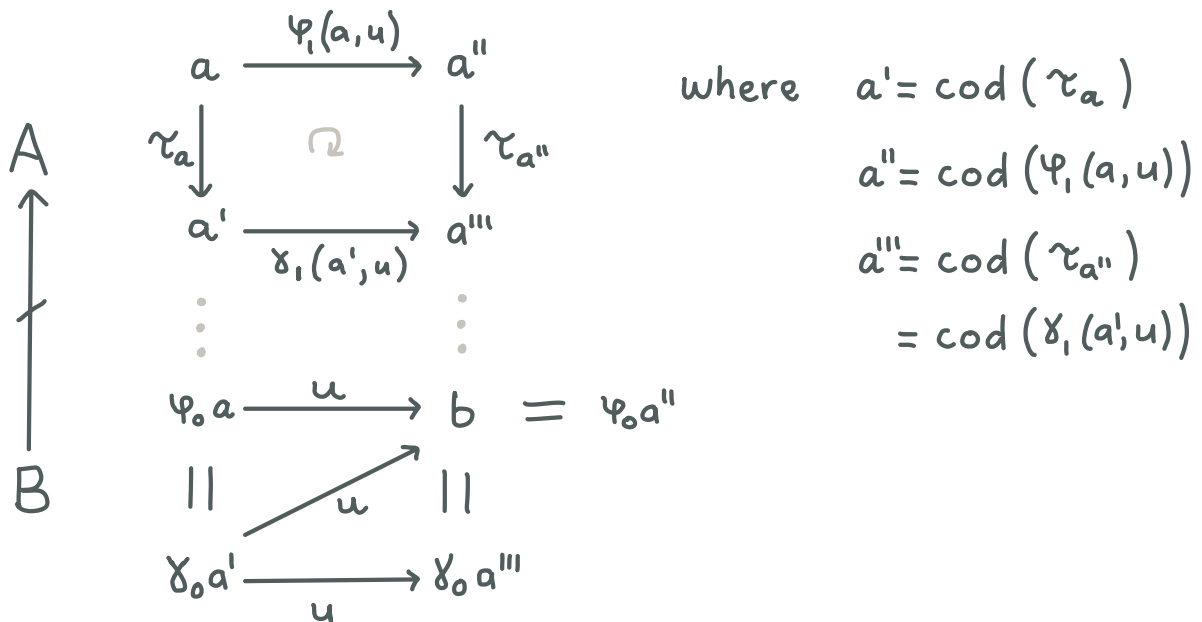
$$\Lambda_1 := A_0 \times_{B_0} B_1$$

□

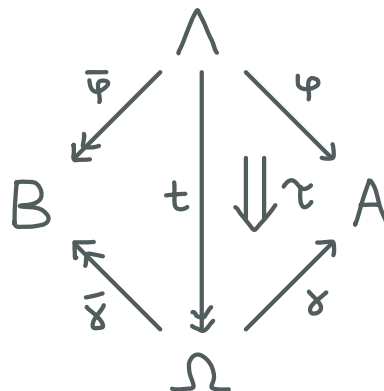
Definition (Aguiar): A **2-cell** between cofunctors,



is an assignment to each object  $a \in A$  of a morphism  $\tau_a: a \rightarrow a'$  in  $A$  such that  $\varphi_0 a = \delta_0 a'$  and for any pair  $(a, u: \varphi_0 a \rightarrow b)$  the following diagram in  $A$  commutes:



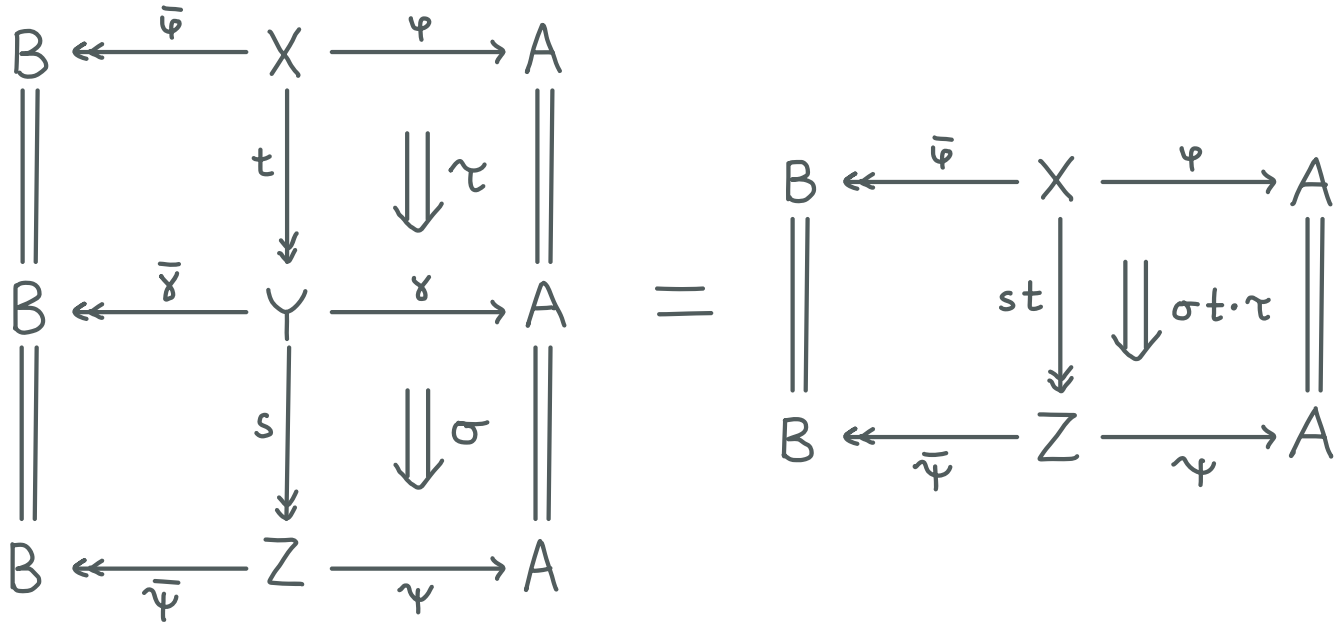
Proposition: Every 2-cell between cofunctors may be represented as a diagram in  $\text{Cat}$ :



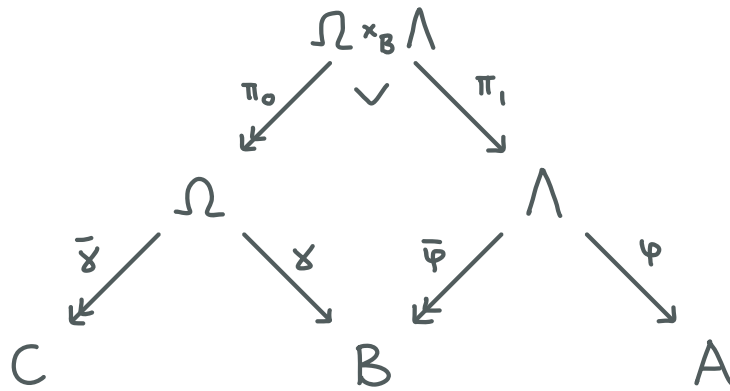
□

Categories, cofunctors and 2-cells form a **2-category**  $\text{Cof}$ .

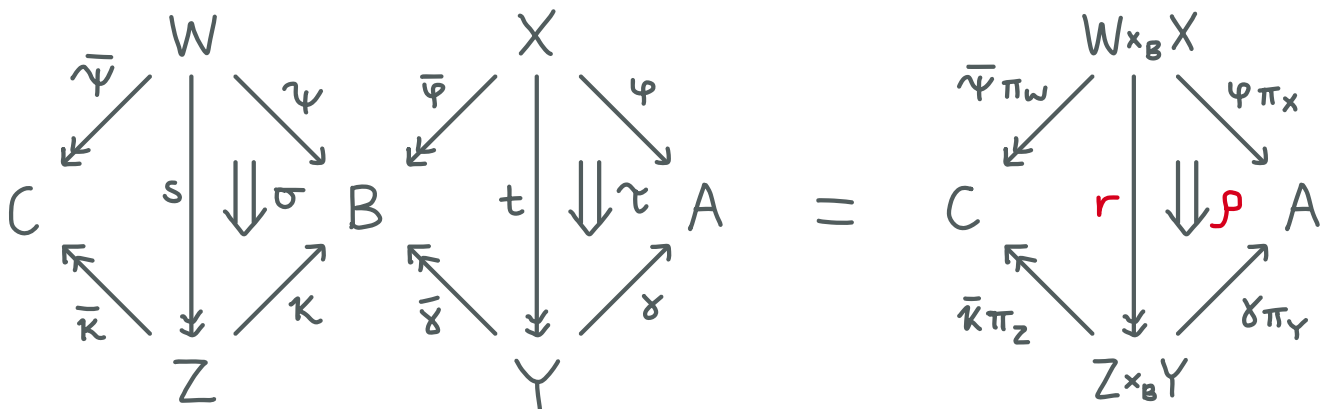
Vertical composition of 2-cells may be understood via pasting in Cat:



Composition of cofunctors corresponds to span composition:



Horizontal composition of 2-cells is more involved.





To obtain the functor  $r$  and natural transformation  $\rho$ , first we must use the universal property of discrete opfibrations:

$$\begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \pi_W \downarrow & & \downarrow \bar{\varphi} \\
 W & \xrightarrow{s} Z \xrightarrow{\kappa} & B
 \end{array}
 \quad \Downarrow \sigma \pi_W \quad = \quad
 \begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \pi_W \downarrow & \Downarrow \theta & \downarrow \bar{\varphi} \\
 W & \xrightarrow{\kappa s} & B
 \end{array}$$

Then the functor  $r$  is constructed by the universal property of the pullback,

$$\begin{array}{ccccc}
 & & \xrightarrow{th} & & \\
 W \times_B X & \xrightarrow{\text{---} r \text{---}} & Z \times_B Y & \xrightarrow{\pi_Y} & Y \\
 \pi_W \downarrow & & \downarrow \pi_Z \lrcorner & & \downarrow \bar{\gamma} \\
 W & \xrightarrow{s} Z \xrightarrow{\kappa} & & & B
 \end{array}$$

while the natural transformation  $\rho$  is the horizontal composite:

$$\begin{array}{ccc}
 W \times_B X & \xrightarrow{\pi_X} & X \\
 \Downarrow \theta & & \downarrow \varphi \\
 W \times_B X & \xrightarrow{\quad} & X \\
 \uparrow h & & \uparrow \gamma t
 \end{array}$$

Note that for  $\mathcal{E}$  with pullbacks, we also have a 2-category  $\text{Cof}(\mathcal{E})$ , where a 2-cell  $\varphi \xRightarrow{\tau} \gamma: B \rightarrow A$  may be specified by:

$$\begin{array}{ccccc}
 & & A_0 & \xrightarrow{\varphi_0} & B_0 \\
 & \swarrow 1 & \downarrow \tau & & \uparrow \gamma_0 \\
 A_0 & \xleftarrow{d_0} & A_1 & \xrightarrow{d_1} & A_0
 \end{array}$$

+ naturality condition

## §2. A right adjoint: the monoid of admissible sections

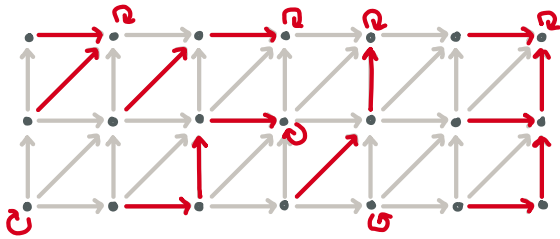
We would like to construct a **right adjoint** to the inclusion.

$$\boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Cof}}$$

How can we construct a monoid from a category?

IDEA: The elements of the monoid are like **vector fields** on the cat.

Example (Spivak, Myers): Consider a full subcategory of  $(\mathbb{N}, \leq) \times (\mathbb{N}, \leq)$ :



At each object, choose an outgoing arrow to get a family:  $(a \xrightarrow{\tau_a} t(a))_{a \in A}$

Definition: The **monoid of admissible sections**  $\Gamma(A)$  for a category  $A$  is a monoid whose:

- elements are families of morphisms  $(a \xrightarrow{\tau_a} t(a))_{a \in A}$
- unit is the family of identities  $(a \xrightarrow{1_a} a)_{a \in A}$
- multiplication is given by:

$$(a \xrightarrow{\tau_a} t(a))_{a \in A} * (a \xrightarrow{\sigma_a} s(a))_{a \in A} = (a \xrightarrow{\sigma_{t(a)} \circ \tau_a} St(a))_{a \in A}$$

We will see there are several equivalent characterisations of this monoid.

Note that  $\Gamma$  defines the object part of the right adjoint:

$$\boxed{\text{Set}} \begin{array}{c} \xrightarrow{F} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Cof}}$$

elements of the right adjoint

$$1 \downarrow \Gamma(A)$$

$$\mathbb{N} \downarrow \Gamma(A)$$

$$\mathbb{N} \downarrow A$$

admissible sections of  $A$

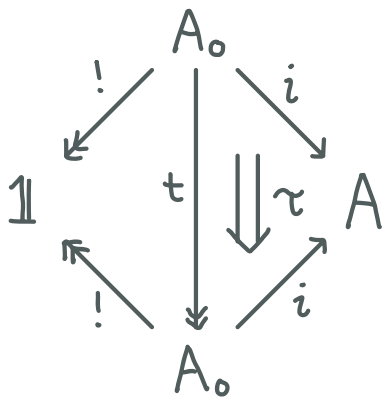
## Examples

- If  $A_0$  is a **discrete category**,  $\Gamma(A_0) \cong \mathbb{1}$
- If  $A$  is a **codiscrete category**,  $\Gamma(A) \cong \text{End}(A_0) = \text{Set}(A_0, A_0)$
- If  $M$  is a **monoid**, then  $\Gamma(M) \cong M$
- For any category  $A$ , the **commutative monoid of endomorphisms** of the identity transformation  $\text{End}(1_A) = \text{Nat}(1_A, 1_A)$  is a submonoid of the centre  $Z(\Gamma(A))$  of  $\Gamma(A)$ .
- (Garner): The **group of extended inner automorphisms** of a category  $A \in \text{Cof}$  is isomorphic to  $\text{Bis}(A)$ , the group of bisections, which is the same as the group of invertible elements of  $\Gamma(A)$ .
- (Aguiar) Recall that a category in  $\text{Cgrp}$  is the same as a crossed module of groups. The monoid of admissible sections of such a category is equivalent to **Whitehead's monoid of derivations** of the crossed module.
- The 2007 paper "External derivations of internal groupoids" by Kasangian, Mantovani, Metere, and Vitale also studies this monoid.
- Every 2-cell  $\tau: \Psi \Rightarrow \chi: B \rightarrow A$  yields an element of  $\Gamma(A)$ .

Proposition: The monoid of admissible sections is isomorphic to the hom category:

$$\Gamma(A) := \text{Cof}(\mathbb{1}, A)$$

PROOF: Consider a morphism in  $\text{Cof}(\mathbb{1}, A)$  given by:



Since  $\mathbb{1}$  is initial in  $\text{Cof}$ , the category  $\text{Cof}(\mathbb{1}, A)$  is a monoid. The components of the natural transformation  $\tau$  give exactly a family of morphisms  $(a \xrightarrow{\tau_a} t(a))_{a \in A}$ , that is, an element of  $\Gamma(A)$ . Moreover, vertical composition of 2-cells corresponds exactly to composition in  $\Gamma(A)$ .  $\square$

Corollary: There is a functor  $\Gamma: \text{Cof} \rightarrow \text{Mon}$ .

PROOF: Given a cofunctor  $B \xrightarrow{\varphi} A$ , we get a morphism of monoids by whiskering  $\text{Cof}(\mathbb{1}, B) \xrightarrow{\varphi} \text{Cof}(\mathbb{1}, A)$ .  $\square$

Note, the statement that  $\Gamma$  is a right adjoint translates to:

$$\text{Cof}(M, A) \cong \text{Mon}(M, \text{Cof}(\mathbb{1}, A))$$

Is this a copower?

Proposition: There is an adjunction of categories:

$$\boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\quad} \\ \perp \\ \xleftarrow{\quad} \\ \Gamma \end{array} \boxed{\text{Cof}}$$

PROOF: The **unit** of the adjunction is a natural isomorphism, and the **counit** can be constructed as the cofunctor  $\Gamma(A) \dashrightarrow A$ :

$$\begin{array}{ccc} A & & x \xrightarrow{\tau_x} t(x) \\ \uparrow \varepsilon_A & & \vdots \qquad \qquad \vdots \\ \Gamma(A) & & \bullet \xrightarrow{\quad} \bullet \\ & & (a \xrightarrow{\tau_a} t(a))_{a \in A} \end{array}$$

In other words, the counit "evaluates" the family at an object  $x \in A$  to get a morphism  $\tau_x \in A$ . One may show the required identities for an adjunction hold.  $\square$

Aside: Recall admissible sections (morphisms in  $\text{Cof}(\mathbb{1}, A)$ ) are the same as cofunctors  $\mathbb{N} \dashrightarrow A$  (objects in  $\text{Cof}(\mathbb{N}, A)$ ). What is the relationship between these categories?

Let  $M$  be a monoid considered as a one-object category. Then we may construct a category  $\hat{M}$  whose objects are morphisms in  $M$  (elements of the monoid), and whose hom sets are given by  $\hat{M}(f, g) = \{m \in M \mid m \cdot f = g \cdot m\}$ . When  $M = \text{Cof}(\mathbb{1}, A)$ , then  $\hat{M} \cong \text{Cof}(\mathbb{N}, A)$ . Does this construction have a universal property?

## Internal perspective

Why is  $\Gamma(A)$  called the monoid of admissible **sections**?

For a category  $\mathcal{E}$  with pullbacks (think  $\mathcal{E} = \text{Set}$ ) we can define a functor  $\Gamma: \text{Cof}(\mathcal{E}) \rightarrow \text{Mon}$  which takes an internal category  $A$  to the monoid  $\Gamma(A)$  whose underlying set is:

$$\Gamma(A) = \{ \sigma : A_0 \rightarrow A_1 \mid d_0 \circ \sigma = 1_{A_0} \} = \{ \text{sections to the domain map } d_0 \}$$

Given a cofunctor  $B \xrightarrow{\varphi} A$ , we can obtain a morphism of monoids  $\Gamma(B) \rightarrow \Gamma(A)$  in the following way:

$$\begin{array}{ccc}
 \Gamma(B) & \xrightarrow{\varphi_0} & \Gamma(\Lambda) & \xrightarrow{\varphi_1} & \Gamma(A) \\
 \begin{array}{ccc} B_0 & \xrightarrow{\sigma} & B_1 \\ \downarrow 1 & & \downarrow d_0 \\ & & B_0 \end{array} & & \begin{array}{ccc} A_0 & \xrightarrow{1 \times \sigma \varphi_0} & \Lambda_1 \\ \downarrow 1 & & \downarrow \pi_0 \\ & & A_0 \end{array} & & \begin{array}{ccc} A_0 & \xrightarrow{\varphi_1(1 \times \sigma \varphi_0)} & A_1 \\ \downarrow 1 & & \downarrow d_0 \\ & & A_0 \end{array}
 \end{array}$$

This is (more or less) how Aguiar originally defined  $\Gamma$ .

However, unlike the previous definition,  $\Gamma$  no longer appears as a right adjoint to the inclusion  $\text{Mon}(\mathcal{E}) \rightarrow \text{Cof}(\mathcal{E})$ .

**Can we fix this?**

Assumption: Let  $\mathcal{E}$  be a category with **finite limits** such that for every object  $B_0 \in \mathcal{E}$ , the functor

$$\mathcal{E} \xrightarrow{B_0 \times (-)} \mathcal{E}/B_0$$

has a **chosen right adjoint**  $\Gamma_{B_0}: \mathcal{E}/B_0 \rightarrow \mathcal{E}$ .

Proposition: Under the assumption above,  $\text{Mon}(\mathcal{E})$  is a coreflective subcategory of  $\text{Cof}(\mathcal{E})$ .

$$\boxed{\text{Mon}(\mathcal{E})} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Cof}(\mathcal{E})}$$

PROOF: Consider an internal category  $B \in \text{Cof}(\mathcal{E})$ . Define:

$$\Gamma(B) := \Gamma_{B_0}(B_1 \xrightarrow{d_0} B_0) \in \mathcal{E}$$

Since  $B_0 \times (-) \dashv \Gamma_{B_0}$ , we have a bijection between:

$$Z \xrightarrow{\hat{\sigma}} \Gamma(B) \quad \iff \quad \begin{array}{ccc} B_0 \times Z & \xrightarrow{\sigma} & B_1 \\ \pi_0 \searrow & & \swarrow d_0 \\ & B_0 & \end{array}$$

In particular, when  $Z = 1$ , we have a correspondence between "elements"  $1 \rightarrow \Gamma(B)$  and sections of  $d_0: B_1 \rightarrow B_0$ .

We can show that  $\Gamma(B)$  is an internal monoid in  $\mathcal{E}$  using the diagrams for  $B$  as an internal category in  $\mathcal{E}$ . For example, the unit of the monoid comes from the identity map:

$$\begin{array}{ccc} B_0 & \xrightarrow{i} & B_1 \\ \downarrow 1 & & \swarrow d_0 \\ & B_0 & \end{array} \quad \iff \quad 1 \xrightarrow{e} \Gamma(B)$$

Multiplication for the monoid (plus the axioms) also follow.

To see **functoriality of  $\Gamma$** , consider a cofunctor  $(\varphi_0, \varphi_1): B \rightarrow A$ .

$$\begin{array}{ccc}
 A_0 \times \Gamma(B) & \xrightarrow{\langle \pi_0, \varepsilon(\varphi_0 \times 1) \rangle} & \Lambda_1 \xrightarrow{\varphi_1} A_1 \\
 \searrow \pi_0 & & \downarrow \pi_0 \quad \swarrow d_0 \\
 & & A_0
 \end{array}
 \rightsquigarrow
 \Gamma(B) \xrightarrow{\Gamma \varphi} \Gamma(A)$$

One may check that the monoid homomorphism axioms hold.

To see that the **unit** of the adjunction is an isomorphism, note that for a monoid  $M$ , we have a bijection:

$$Z \longrightarrow \Gamma(M) \rightsquigarrow Z \begin{array}{c} \xrightarrow{\quad} M \\ \swarrow ! \\ \downarrow ! \\ 1 \end{array}$$

We can easily show that  $M \cong \Gamma(M)$  follows.

Finally the **counit** of the adjunction at an internal category  $B$  is the internal cofunctor (given as a span of internal functors) below:

$$\begin{array}{ccccc}
 1 & \xleftarrow{!} & B_0 & \xrightarrow{1} & B_0 \\
 \uparrow & & \uparrow \pi_0 & & \uparrow d_0 \\
 \Gamma(B) & \xleftarrow{\pi_1} & B_0 \times \Gamma(B) & \xrightarrow{\varepsilon} & B_1 \\
 \downarrow & & \downarrow d_1 \varepsilon & & \downarrow d_1 \\
 1 & \xleftarrow{!} & B_0 & \xrightarrow{1} & B_0
 \end{array}$$

$\varepsilon$  is the counit of the adjunction  $B_0 \times (-) \dashv \Gamma_{B_0}$

□



## SUMMARY (IF SHORT ON TIME)

- In §1 we saw that there is a 2-category  $\mathcal{Cof}$  of small categories, **cofunctors**, and 2-cells.
- Furthermore, we saw that cofunctors and their 2-cells could be understood via **diagrams in  $\mathcal{Cat}$** .
- In §2 we proved that the category of monoids is a **coreflective subcategory** of  $\mathcal{Cof}$ :

$$\boxed{\text{Mon}} \begin{array}{c} \xrightarrow{\text{f.f.}} \\ \perp \\ \xleftarrow{\Gamma} \end{array} \boxed{\text{Cof}}$$

- We saw that the right adjoint, which takes a category to its **monoid of admissible sections**, has several characterisations including as the hom category  $\Gamma(-) = \mathcal{Cof}(\mathbb{1}, -)$ .
- Moreover, this adjunction generalises to the **internal setting** with suitable assumptions on  $\mathcal{E}$ .

## DIRECTIONS FOR FUTURE WORK

- It is possible to construct a **double category** of functors and cofunctors from  $\mathcal{Span}$  using Pare's **retrocells**. What can we learn from this perspective? What properties of  $\mathcal{Span}$  allow us to represent cofunctors as spans?
- Replacing  $\mathcal{E}$  with a suitable **monoidal category  $\mathcal{V}$** , we can generalise the above adjunction to involve monoids and categories in  $\mathcal{V}$  (this is what Aguiar did) and obtain string diagram proofs.
- Are there **other examples** of  $\Gamma(A)$  in the literature? What interesting information about  $A$  does it contain? Applications?